

## 16. Seeded Lossless Rank Extractors

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Let  $\mathbb{F} = \overline{\mathbb{F}}$  (e.g.  $\mathbb{F} = \mathbb{C}$ ). Let  $M = (M_i)_{i \in I}$  be a finite collection of matrices  $M_i \in \mathbb{F}^{m \times n}$ , where  $m \leq n$ .

Def We say  $M$  is an  $(m', \epsilon)$ - (seeded) lossless rank condenser if for every  $A \in \mathbb{F}^{n \times m'}$  of rank  $m'$ ,

$$\Pr_{i \in I} [\text{rank}(M_i \cdot A) = m'] \geq 1 - \epsilon.$$

If  $m = m'$ , we say  $M$  is an  $\epsilon$ - (seeded) lossless rank extractor.

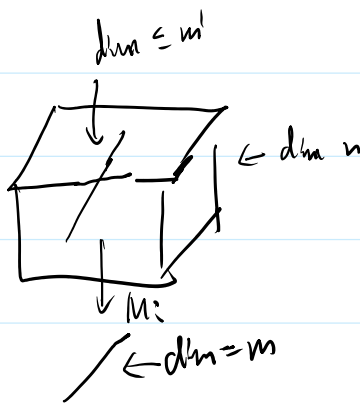
Let  $V \subseteq \mathbb{F}^n$  be the column span of  $A$ . Then  $\dim V = \text{rank}(A) = m'$

View  $M_i$  as a linear map  $\mathbb{F}^n \rightarrow \mathbb{F}^m$ . Then  $\text{rank}(M_i \cdot A) = m' \Leftrightarrow \dim(M_i(V)) = m'$ .

So  $A$  is an  $(m', \epsilon)$ - lossless rank condenser iff all but  $\leq \epsilon$ -fraction of  $M_i$  preserve the dimension of  $V$  (i.e.  $\dim V = \dim M_i(V)$ ) for all linear subspace

$V \subseteq \mathbb{F}^n$  of dimension (at most)  $m'$ .

$\Rightarrow$  If  $M_i$  is injective on  $V$  of  $\dim m'$ , then it is also injective on subspaces of  $V$ .



"seeded" means we consider  $M_i$  chosen from a family  $M$ , not a single matrix  $M_i$ .

A single matrix can't work (why?)

"lossless" means  $\text{rank}(A) = \text{rank}(M_i \cdot A)$ , or equivalently,  $\dim V = \dim(M_i(V))$ .

A construction of  $\epsilon$ -lossless rank extractors. (Forbes - Shpilka '12)

finite  $\rightarrow$  Let  $I \subseteq \mathbb{F}[t]$ . Let  $\omega \in \mathbb{F}$  s.t. the multiplicative order of  $\omega$  is at least  $n$ , i.e.,  $1, \omega, \dots, \omega^{n-1}$  are distinct.

For  $a \in I$ , let  $M_a = (\omega^{it} a)^{j-1}_{1 \leq i \leq m, 1 \leq t < n} \in \mathbb{F}^{m \times n}$ .

finite

For  $a \in I$ , let  $M_a = (\omega^{i \cdot j} a)^{j-1}_{i \leq i \leq m, 1 \leq j \leq n} \in \mathbb{F}^{m \times n}$ .

Thm (Forbes-Shyilkal'12, Forbes-Saptharishi-Shyilkal'13)

$M = (M_a)_{a \in I}$  is an  $\varepsilon$ -lossless rank extractor with  $\varepsilon = \frac{m(n-m)}{|I|}$ .

i.e. for every  $A \in \mathbb{F}^{m \times n}$ , the number of  $a \in I$  s.t.  $\text{rank}(M_a \cdot A) < m$  is at most  $m(n-m)$ .

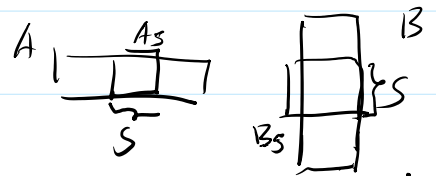
To prove the theorem, we need:

Lemma (Cauchy-Binet) Let  $A \in \mathbb{F}^{m \times n}$  and  $B \in \mathbb{F}^{n \times m}$ , where  $m \leq n$ . Then

$$\det(A \cdot B) = \sum_{\substack{S \subseteq [n] \\ |S|=m}} \det(A_S) \cdot \det(B_S)$$

← i.e.  $S$  ranges over  $m$ -subsets of  $[n] = \{1, \dots, n\}$ .

where  $A_S$  is the  $m \times m$  submatrix with cols selected by  $S$ ,  
and  $B_S$  is the  $m \times m$  submatrix with rows selected by  $S$ .



Pf of the Lemma: Let  $a_1, \dots, a_m$  be the rows of  $A$ .  
&  $b_1, \dots, b_m$  the columns of  $B$ .

$$\text{Note } A \cdot B = \begin{pmatrix} a_1 \cdot B \\ \vdots \\ a_m \cdot B \end{pmatrix} = \begin{pmatrix} A \cdot b_1 & \dots & A \cdot b_m \end{pmatrix}.$$

As  $\det$  is multilinear in its rows and columns,  $\det(A \cdot B)$  is multilinear in  $a_1, \dots, a_m$  and  $b_1, \dots, b_m$  (over  $\mathbb{F}$ ).

i.e. let  $f(a_1, \dots, a_m, b_1, \dots, b_m) = \det(A \cdot B)$ .

Then  $f(a_1, \dots, r a_i + s a'_i, \dots, b_1, \dots, b_m) = r f(\dots a_i, \dots) + s f(\dots a'_i, \dots)$   
and similarly for  $b_i$ 's.  $r, s \in \mathbb{F}$ .

Write each  $a_i$  and  $b_i$  as a linear combination of the standard basis  $e_1, \dots, e_n$ .

By multilinearity, we may assume  $a_i, b_i \in \{e_1, \dots, e_n\}$ .

Say  $a_i = e_{j_i}, b_i = e_{j'_i}, S_A = \{j_1, \dots, j_m\}, S_B = \{j'_1, \dots, j'_m\}$ .

If  $|S_A| < m$  or  $|S_B| < m$ ,  $\det(A \cdot B) = 0$  and  $\sum_S \det(A_S) \cdot \det(B_S) = 0$ .

If  $S_A \neq S_B$ , again  $\det(A \cdot B) = 0$  and  $\sum_S \det(A_S) \cdot \det(B_S) = 0$ .

If  $S_A \neq S_B$ , again  $\det(A|B) = 0$  and  $\sum \det(A_i) \cdot \det(B_i) = 0$ .  
 So we may assume  $S_A = S_B$  and  $|S_A| = |S_B| = m$ .

Then  $A \cdot B = A_{S_A} \cdot B_{S_B}$ . So the claim holds.  
 zero outside  $S_A$     zero outside  $S_B$ .

□ substitute  $X$  by  $a$  within each entry.

Proof of Thm: Consider  $M(x) = (w^{i-1} x)^{j-1}_{1 \leq i \leq m, 1 \leq j \leq n} \in \mathbb{F}[x]^{m \times n}$ . So  $M_a = M(a)$ .

Consider  $A = (a_{ij})_{1 \leq i \leq n, 1 \leq j \leq m} \in \mathbb{F}^{n \times m}$ . Let  $P(x) = \det(M(x) \cdot A) \in \mathbb{F}[x]$ .

Claim:  $P(x) \neq 0$ ,  $\deg(P) \leq \sum_{i=n-m+1}^n (i-1)$ , and  $X^{\sum_{i=1}^m (i-1)}$  divides  $P(x)$ .

Pf of the claim: By Cauchy-Binet,

$$\begin{aligned} |P(x)| &= \sum_S \det((M(x))_S) \cdot \det(A_S) \\ &= \sum_{\substack{1 \leq d_1 < \dots < d_m \leq n \\ S = \{d_1, \dots, d_m\}}} \det((w^{i-1} x)^{d_j-1}_{1 \leq i \leq m, 1 \leq j \leq m}) \cdot \det(A_S) \\ &= \sum_{1 \leq d_1 < \dots < d_m \leq n} \det((w^{i-1})^{d_j-1}) \cdot \det(A_S) \cdot X^{\sum_{j=1}^m (d_j-1)} \quad (*) \end{aligned}$$

choose  $1 \leq d_1 < \dots < d_m \leq n$  s.t.  $\det(A_S) \neq 0$  and  $\sum_{j=1}^m (d_j-1)$  is minimized.

By exchangeability,  $(d_1, \dots, d_m)$  is unique. The corresponding  $\det((w^{i-1})^{d_j-1}) = \det((w^{d_j-1})^{i-1}) \neq 0$   
 $\nearrow$  Vandermonde matrix  
 $1, w, \dots, w^{n-1}$  are distinct.

So  $P(x) \neq 0$ ,  $\deg(P) \leq \sum_{i=n-m+1}^n (i-1)$ , and  $X^{\sum_{i=1}^m (i-1)}$  divides  $P(x)$  by (\*).

This proves the claim.

$\Rightarrow$   $P(x)$  has at most  $\left(\sum_{i=n-m+1}^n (i-1)\right) - \left(\sum_{i=1}^m (i-1)\right) = m(n-m)$  roots in  $\mathbb{F} \setminus \{0\} \geq I$ .

$$\langle \mathcal{D} \mid \text{rank}(M \cdot A) = m \rangle = \mathcal{D} \cdot \left[ \det(M(x) \cdot A) \neq 0 \right]$$

$$\begin{aligned} \text{So } \Pr_{a \in I} [\text{rank}(M_a A) = m] &= \Pr_{a \in I} [\det(M(x) \cdot A)|_{x=a} \neq 0] \\ &= \Pr_{a \in I} [P(a) \neq 0] \leq \frac{m(n-m)}{|I|}. \quad \square \end{aligned}$$

By the theorem, if  $|I| > m(n-m)$ , then for every  $A \in \mathbb{F}^{n \times m}$ ,  $\exists M_a \in \mathcal{M}$  s.t.  
 $\text{rank}(M_a A) = \text{rank}(A)$ .

This is tight:

Thm: If  $|I| \leq m(n-m)$ , for any collection  $\mathcal{M} = (M_a)_{a \in I}$  of matrices  $M_a \in \mathbb{F}^{m \times n}$ ,  
 $\exists A \in \mathbb{F}^{n \times m}$  of rank  $m$  s.t.  $\text{rank}(M_a A) < m$  for all  $a \in I$

Pf sketch: The set of  $m$ -dimensional subspaces of  $\mathbb{F}^n$   
 is called the Grassmannian  $\text{Gr}_{\mathbb{F}}(m, n) = \text{Gr}(m, n)$

It embeds in the projective space  $\mathbb{P} \binom{n}{m}$ :



$V \mapsto (A_s)_{s \in \binom{[n]}{m}}$ , where  $A \in \mathbb{F}^{n \times m}$  s.t.  $V$  is the column space of  $A$ .

For each  $a \in I$ ,  $\text{rank}(M_a A) < m \Leftrightarrow \det(M_a A) = 0 \Leftrightarrow \sum_s \det(M_a)_s \det(A_s) = 0$   
 (Cayley-Binet)

This gives a linear constraint  $C_a$  in the coordinates  $A_s$ .

As  $\dim \text{Gr}(m, n) = m(n-m) > |I|$ , there exists  $V \in \text{Gr}(m, n)$   
 satisfying  $C_a$  for all  $a \in I$ , i.e.,  $\text{rank}(M_a A) < m$  for all  $a \in I$ .  $\square$